# **CSE113: Parallel Programming**

- Topics:
  - Example Questions
  - Processes







## Announcements

- HW 4 grades will be released this week.
- HW 5 is due today.
- SETs are out, please do them! It helps us out a lot.

How many API calls do Barrier objects have?

○ 0
 ○ 1
 ○ 2
 ○ 3

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A barrier call emits which of the following events? Check all that apply

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TrueFalse

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TrueFalse

If the mutex is protecting concurrent writes, there is non-determinism.

Write a few sentences about what you think the best interface for parallel programming is, that is, do you think it is Atomics? Mutexes? Concurrent Data Structures? Barriers? Or even maybe the compiler should simply do it all automatically? Or is it some combination of the above? What are the trade-offs involved?

# Sample Questions

Separate file

## Zombies

### ∎ldea

When process terminates, still consumes system resources

Various tables maintained by OS

Called a "zombie"

Living corpse, half alive and half dead

### ■Reaping

Performed by parent on terminated child (using wait or waitpid)

Parent is given exit status information

Kernel discards process

### What if parent doesn't reap?

If any parent terminates without reaping a child, then child will be reaped by init process (pid == 1)

So, only need explicit reaping in long-running processes

•e.g., shells and servers

Zombie Example	<pre>void fork7( {     if (for         /*         pri         exi     } else     pri </pre>	<pre>) k() == 0) { Child */ intf("Terminating Child, PID = %d\n",     getpid()); it(0); { intf("Running Parent, PID = %d\n",</pre>	
<pre>linux&gt; ./forks 7 &amp; [1] 6639 Running Parent, PID = 6639 Terminating Child, PID = 6640 linux&gt; pa</pre>	- whi } }	<pre>getpid()); while (1) ; /* Infinite loop */ }</pre>	
Finux> ps         PID TTY       TIME CMD         6585 ttyp9       00:00:00 tcsh         6639 ttyp9       00:00:03 forks         6640 ttyp9       00:00:00 forks         6641 ttyp9       00:00:00 ps         linux> kill 6639         [1]       Terminated         linux> ps         PID TTY       TIME CMD         6585 ttyp9       00:00:00 tcsh         6585 ttyp9       00:00:00 tcsh	<defunct></defunct>	<ul> <li><b>ps</b> shows child process as "defunct"</li> <li>Killing parent allows child to be reaped by <b>init</b></li> </ul>	

## **Orphan process:** Nonterminating **Child process**

{

```
linux> ./forks 8
Terminating Parent, PID = 6675
Running Child, PID = 6676
linux> ps
 PID TTY
                  TIME CMD
 6585 ttyp9 00:00:00 tcsh
 6676 ttyp9 00:00:06 forks
 6677 ttyp9 00:00:00 ps
linux> kill 6676
linux> ps
 PID TTY
                  TIME CMD
 6585 ttyp9
             00:00:00 tcsh
 6678 ttyp9
              00:00:00 ps
```

#### void fork8() if (fork() == 0) { /\* Child \*/ printf("Running Child, PID = $d\n$ ", getpid()); while (1) ; /\* Infinite loop \*/ } else { printf("Terminating Parent, PID = %d\n", getpid()); exit(0); }

■Child process still active even though parent has terminated. The process init adopts the process. Daemons can be created this way.

■Must kill explicitly, or else will keep running indefinitely

### wait: Synchronizing with Children

#### Parent reaps child by calling the wait function

#### int wait(int \*child\_status)

suspends current process until one of its children terminates
return value is the pid of the child process that terminated
if child\_status != NULL, then the object it points to will be set to a status indicating why the child process terminated

### wait: Synchronizing with Children

```
void fork9() {
    int child_status;

    if (fork() == 0) {
        printf("HC: hello from child\n");
    }
    else {
        printf("HP: hello from parent\n");
        wait(&child_status);
        printf("CT: child has terminated\n");
    }
    printf("Bye\n");
    exit();
}
```



## wait() Example

If multiple children completed, will take in arbitrary order
 Can use macros WIFEXITED and WEXITSTATUS to get information about exit status (W for wait)

```
void fork10()
Ł
   pid t pid[N];
   int i;
    int child status;
    for (i = 0; i < N; i++)
       if ((pid[i] = fork()) == 0)
           exit(100+i); /* Child */
    for (i = 0; i < N; i++) {
       pid t wpid = wait(&child status);
       if (WIFEXITED(child status))
           printf("Child %d terminated with exit status %d\n",
                  wpid, WEXITSTATUS(child status));
       else
           printf("Child %d terminate abnormally\n", wpid);
```

### waitpid(): Waiting for a Specific Process

```
waitpid(pid, &status, options)
```

suspends current process until specific process terminates
various options

```
void fork11()
{
   pid t pid[N];
    int i;
    int child status;
    for (i = 0; i < N; i++)
       if ((pid[i] = fork()) == 0)
           exit(100+i); /* Child */
    for (i = N-1; i \ge 0; i--) {
       pid t wpid = waitpid(pid[i], &child status, 0);
       if (WIFEXITED(child status))
           printf("Child %d terminated with exit status %d\n",
                  wpid, WEXITSTATUS(child status));
       else
           printf("Child %d terminated abnormally\n", wpid);
```

### execve: Loading and Running Programs

```
int execve(
    char *filename,
    char *argv[],
    char *envp[]
```

```
)
```

#### **Loads and runs in current process:**

- ■Executable filename
- With argument list argv
- And environment variable list envp
- Does not return (unless error)

#### **Overwrites code, data, and stack**

keeps pid, open files

### Environment variables:

- "name=value" strings
- ■Use functions getenv and putenv to
  access environment variables.

The v and e comes from the fact that it takes an argument argv, envp to the vector of arguments and environment variables to the program

envp[n] == NULL	S
envp[n-1]	
envp[0]	
argv[argc] == NULL	
argv[argc-1]	
argv[0]	
envp	
argv	
argc	
Stack frame for	
main	Ģ

Stack bottom

Stack top

### execve Example

